

Winter 2010 SYSC 5201 Midterm

90 Minutes

Close book

1. (30%) Provide simple and straight answers to the following questions.
 - a. What is the main challenge in network edge?
Various access technologies and limited bandwidth
 - b. What is the main disadvantage of radio media?
Interference
 - c. What is Little's Law?
 $L = \lambda T$
 - d. What are the main differences between client and server?
Server is always on while client is on intermittently. Client initiates service.
 - e. Which factors affect the transmission delay of a packet?
Packet size and bandwidth
 - f. Which factors affect propagation delay?
Distance and signal propagation speed.
 - g. What is the main difference between persistent and non-persistent HTTP?
Persistent HTTP uses one TCP connection while non-persistent HTTP uses multiple.
 - h. How does a distributed peer-to-peer architecture keep and access index information?
Using DHT
 - i. What is the main difference between SMTP and POP3?
SMTP is a push protocol while POP3 is a pull protocol.
 - j. What is a socket?
An interface and API between application layer and transport layer.

2. (30%) Consider a network with three packet switches. Packets arrive at switches 1, 2, 3 in accordance with Poisson processes having respective rates 5, 10, 15. The transmission times at the three switches are exponential with respective rates 10, 50, 100. A packet completing service at switch 1 is equally likely to either (a) go to switch 2, (b) go to switch 3, or (c) leave the system. A packet departing at switch 2 always go to switch 3. A departure packet from switch 3 is equally likely to either go to switch 2 or leave the system. Ignore propagation and processing delays.
 - a. What is the average number of packets in the network (consisting of all three switches)?
 - b. What is the average time a packet spends in the network?
 - c. What is the probability that there are more than 4 packets in switch 1?

Solution:

- a) 82/13
- b) 41/195

$$c) \sum_{n=5}^{\infty} \left(\frac{1}{2}\right)^n \left(\frac{1}{2}\right) = \left(\frac{1}{2}\right)^6 \left(\sum_{n=0}^{\infty} \left(\frac{1}{2}\right)^n\right) = \frac{1}{32}$$

3. (40%) Consider distributing a file of F bits to N peers using a P2P architecture. Assume a fluid model. For simplicity assume that d_{min} is very large, so that peer download bandwidth is never a bottleneck. Suppose that $u_s \leq (u_s + u_1 + \dots + u_n)/N$. Specify a distribution scheme that has a distribution time of F/u_s .

Solution:

Define $u = u_1 + u_2 + \dots + u_N$. By assumption

$$u_s \leq (u_s + u)/N \quad \text{Equation 1}$$

Divide the file into N parts, with the i^{th} part having size $(u_i/u)F$. The server transmits the i^{th} part to peer i at rate $r_i = (u_i/u)u_s$. Note that $r_1 + r_2 + \dots + r_N = u_s$, so that the aggregate server rate does not exceed the link rate of the server. Also have each peer i forward the bits it receives to each of the $N-1$ peers at rate r_i . The aggregate forwarding rate by peer i is $(N-1)r_i$. We have

$$(N-1)r_i = (N-1)(u_s u_i)/u \leq u_i,$$

where the last inequality follows from Equation 1. Thus the aggregate forwarding rate of peer i is less than its link rate u_i .

In this distribution scheme, peer i receives bits at an aggregate rate of

$$r_i + \sum_{j < i} r_j = u_s$$

Thus each peer receives the file in F/u_s .